

# A probabilistic model for asynchronous iterative methods

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 $\begin{array}{l} \textit{Outline} \cdot \mathsf{Motivation} \cdot \mathsf{Model} \cdot \mathsf{Analysis} \cdot \mathsf{Summary} \\ \textbf{Outline} \end{array}$ 

- Motivation
- Asynchronous Richardson: Model and analysis
- Asynchronous Schwarz: Model and analysis
- Summary and outlook



 $Outline \cdot \textit{Motivation} \cdot Model \cdot Analysis \cdot Summary$ 

# Motivation

• Compute systems are increasingly heterogeneous





One Frontier node (#1 in top500)

- Compute systems are increasingly heterogeneous
- Compute systems are increasingly hierarchical



NVIDIA GPU schematic



 $\begin{array}{l} \text{Outline} \cdot \textit{Motivation} \cdot \text{Model} \cdot \text{Analysis} \cdot \text{Summary} \\ \text{Motivation} \end{array}$ 

- Compute systems are increasingly heterogeneous
- Compute systems are increasingly hierarchical
- Synchronization bottlenecks restrict scalability





Outline  $\cdot$  *Motivation*  $\cdot$  Model  $\cdot$  Analysis  $\cdot$  Summary Motivation

- Compute systems are increasingly heterogeneous
- Compute systems are increasingly hierarchical
- Synchronization bottlenecks restrict scalability
- Fault tolerance can improve robustness





# A typical computational physics workflow



Domain (left<sup>1</sup>) and simulation (right<sup>2</sup>)

<sup>1</sup>Liu et. al, Oct, 2018 <sup>2</sup><u>https://amrex-combustion.github.io/</u>

**3** 2024.05.27 Pratik Nayak - Probabilistic modeling for asynchronous methods





# A typical computational physics workflow





# A typical computational physics workflow



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Outline · *Motivation* · Model · Analysis · Summary Solving linear systems Direct methods

- Provide solution in a fixed number of steps to required precision
- Computationally more expensive ~  $\mathcal{O}(n^3)$





#### Iterative methods

- Successively approximate solution
- Computationally cheaper, dependent on matrix properties ~  $\mathcal{O}(kn^2), \;\; k << n$
- Efficient when exact solution not necessary





# Conjugate gradient: Operation breakdown





# Synchronization-free methods

- Don't explicitly synchronize
- Take latest available data from your neighbor
- Compute on the partially updated data
- Convergence guaranteed only under certain conditions



GPU 0	GPU 1	GPU 2
GPU 3	GPU 4	GPU 5
GPU 6	GPU 7	GPU 8



#### Asynchronous methods: Convergence conditions<sup>1</sup>

The asynchronous iteration is defined as 
$$x_{k+1}^l = \begin{cases} x_k^l & \text{if } l \notin \mathcal{G}_k \\ f_l(x_{s_1(k)}^1, \cdots, x_{s_n(k)}^n) & \text{if } l \in \mathcal{G}_k \end{cases}$$

With iteration subsets,  $\mathcal{G}_k \subseteq \{1, ..., J\}$  and delay subsets  $\mathcal{S}_k = s_1(k), \cdots, s_n(k)$ , the conditions necessary for convergence are

the sets  $\{k \mid j \in \mathcal{G}_k\}$  are unbounded for j = 1, ..., J

$$s_j(k) \le k-1$$
  $\forall \ j,k$  Only previous iterations can be used

$$lim_{k\to\infty}s_j(k) = \infty$$
 for  $j = 1, ..., J$  Use the latest local update

Each component needs to be continuously updated

The asynchronous iteration is denoted by  $(\mathcal{F}, x_0, \mathcal{G}, \mathcal{S})$ <sup>1</sup>Chazan and Miranker, Chaotic relaxation



# Asynchronous methods: Probabilistic model

The asynchronous iteration is defined by  $(\mathcal{F}, x_0, \mathcal{G}, \mathcal{S})$ , as defined before can be modeled with a probabilistic model. Let  $k_{\mathcal{F}}$  be the latest local update and  $\mathcal{D} \sim \mathcal{P}(x)$  be some probability distribution where the delays are sampled from.  $\mathcal{S}_{\mathcal{P}} = k_{\mathcal{F}} - \mathcal{D}$ , therefore gives us iteration from which we incorporate information. The probabilistic model samples the delays,  $p(k) \in \mathcal{D} \sim \mathcal{P}$ , subject to the conditions,

- 1. Positivity:  $p(k) \ge 0$
- 2. Validity:  $p(k) \leq k 1$ , delays cannot be greater than the current iteration.
- 3. Using latest available update: p(k+1) p(k) > 0
- 4. No stagnation:  $p(k+1) \rightarrow \infty$
- 5. Offset:

$$p(k) egin{cases} = 0, & \textit{if} \ \ k < o_{\mathcal{F}} \ \in \mathscr{P}_S, & \textit{otherwise} \end{cases}$$

where  $o_{\mathcal{F}} \geq 0$  is the iteration offset at which the delays are introduced.



# Asynchronous methods: Probability distributions

- Aim to model realistic systems
- Exponential: High probability of no delays
- Normal: Delays distributed around the mean.





# Asynchronous methods: Empirical estimation

- Linear fit on the residual, and convergence rate given by the slope of the fit.
- Assuming linear convergence, number of iterations for convergence

 $i = \frac{\log(\tau_{final})}{\log(\varrho(A))}$ 

• Therefore,  $0 < \rho(A) \le 1$ 



 $Outline \ \cdot \ Motivation \ \cdot \ \textit{Model} \ \cdot \ Analysis \ \cdot \ Summary$ 



# Asynchronous Richardson: Problem

- Laplace 2D problem
- Symmetric and positive definite matrix
- On a 5 x 5 grid, 25 degrees of freedom





### Asynchronous Richardson: Random sampling algorithm

- 1:  $A, b, x, \omega$
- 2:  $r \leftarrow b Ax, e \leftarrow \mathbf{0}$
- 3: for  $k < N_{iter}$  do
- 4:  $d \leftarrow \texttt{RANDOM}_\texttt{SAMPLE}(k, a, b)$
- 5:  $r \leftarrow b Ax(d)$
- 6: if  $\|r\| < \tau$  then
- 7: break
- 8: end if
- 9:  $e \leftarrow \texttt{PRECOND}(r)$
- 10:  $x \leftarrow x + \omega e$
- 11: end for



#### Asynchronous Richardson: Exponential distribution







### Asynchronous Richardson: Half-normal distribution



 $\omega=\,0.8.$ 



#### Asynchronous Richardson: Relaxation parameter





# Asynchronous Schwarz: Definition<sup>1</sup>

The asynchronous iteration is defined as  $x_{k+1}^l = \begin{cases} x_k^l & \text{if } l \notin \mathcal{G}_k \\ \sum_{j=1}^J D_{l,j}^{(k)} & y_{s_j(k)}^j & \text{if } l \in \mathcal{G}_k \end{cases}$ With iteration subsets,  $\mathcal{G}_k \subseteq \{1, ..., J\}$  and delay subsets  $\mathcal{S}_k = s_1(k), \cdots, s_n(k)$ , and  $y_{s_i(k)}^j = M_j^{-1}(N_j x_{s_i(k)}^j + f)$ 

With the splittings  $A = M_j - N_j$   $j = 1, \dots, J$ , and the weighting matrices  $D_{l,j}^{(k)}$  such that  $\sum_{j=1}^{J} D_{l,j}^{(k)} = I \quad \forall l, k$ 

The asynchronous iteration is denoted by  $(\mathcal{F}_{schw}, x_0, \mathcal{G}, \mathcal{S})$ ; called the asynchronous Schwarz method.

<sup>1</sup>Frommer and Szyld



# Asynchronous Two-level Schwarz: Definition<sup>1</sup>

A coarse level enables robustness for domain decomposition methods, enhancing the exchange of global information. We emulate the action of a global coarse matrix,  $A_0^{-1}$ , and incorporate it into the solution as shown below

$$\mathbf{u}^{k+1/2} = \mathbf{u}^{k} + \sum_{i=1}^{p} R_{i}^{\top} D_{i} A_{i}^{-1} R_{i} (\mathbf{f} - A \mathbf{u}^{k})$$
$$\mathbf{u}^{k+1} = \mathbf{u}^{k+1/2} + R_{0}^{\top} A_{0}^{-1} R_{0} (\mathbf{f} - A \mathbf{u}^{k+1/2})$$

Coarse correction

<sup>1</sup>Szyld et.al

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# Asynchronous Richardson: Problem

- Laplace 2D problem
- Symmetric and positive definite matrix
- On a 8 x 8 grid, 64 degrees of freedom





#### Asynchronous Schwarz: Two-level probabilistic model

1:  $A, b, x, \omega$ 2:  $r \leftarrow b - Ax, e \leftarrow 0$ 3: for  $k < N_{iter}$  do if  $\|r\| < \tau$  then 4: break 5: 6: end if for  $j < N_{domains}$  do 7:  $d_i \leftarrow s_i(k) \in \mathcal{S}_k$ 8: ▷ Get the sampled delay  $r_i \leftarrow b_i - A_i x_i^{d_j}$ 9: > Communicate and update global solution with delay  $e_j = A_{j,j} ackslash r_j \ x_j^{k+1} := x_j^k + e_j$ 10: ▷ Solve locally 11: ▷ Update local solution 12: end for if coarse corr == additive then 13: if  $mod(f, N_{iter}) == 0$  then 14: > Update only for coarse correction frequency, /  $e_{coarse} = A_{coarse} \backslash Rr$ 15: Compute coarse correction  $x^{k+1} := x^{k+1} + \alpha_{coarse} Pe_{coarse}$ 16: > Weighted additive coarse correction end if 17: else if coarse\_corr == multiplicative then 18:  $r \leftarrow b - Ax^{k+1}$ 19: Compute updated residual  $e_{coarse} = A_{coarse} \backslash Rr$ 20: Compute coarse correction  $x^{k+1} := x^{k+1} + Pe_{coarse}$ 21: ▷ Multiplicative coarse correction end if 22: 23: end for



#### Asynchronous Schwarz: Probabilistic model, coarse correct





#### Asynchronous Schwarz: Probabilistic model, delay



21





# Asynchronous Schwarz: Probabilistic model, coarse freq







# Asynchronous Schwarz: Probabilistic model, coarse levels





1 coarse level





# Summary and future work

- Asynchronous methods can be modeled with a probabilistic approach
- Effects of delays can be controlled to balance between convergence and robustness
- Extending these methods to faster converging methods such as Krylov methods can be very interesting.
- Utilizing probabilistic linear solvers<sup>1</sup> within this framework could allow to bundle in uncertainty within the solution and the system instabilities.

<sup>1</sup>Hennig et.al, Probabilistic linear solvers



# Asynchronous methods: Richardson convergence<sup>1</sup>

The asynchronous iteration is defined by  $(\mathcal{L}, x_0, \mathcal{G}, \mathcal{S})$  with an iteration matrix T, converges if

(a) converges if there exists a positive vector, v and a scalar  $\alpha < 1$  such that  $|T|v \leq \alpha v$ .

(b) The spectral radius of the iteration matrix, T, ρ(|T|) < 1.</li>
Similarly the Richardson iteration with (T<sub>ω</sub>, x<sub>0</sub>, 𝔅, 𝔅), T<sub>ω</sub> = I − ωM<sup>-1</sup>A, A = M − N, M = A<sub>ii</sub> and N = A<sub>i≠j</sub> converges if
(a) ρ(|T<sub>ω</sub>|) = α < 1</li>
(b) 0 < ω < 2/(1+α)</li>

<sup>1</sup>Chazan and Miranker, Chaotic relaxation